



Proceedings

THIRTY-SEVENTH ANNUAL ALLERTON CONFERENCE ON COMMUNICATION, CONTROL AND COMPUTING

September 22 - 24, 1999

Allerton House, Monticello, Illinois

Sponsored by the

Coordinated Science Laboratory and the

Department of Electrical and Computer Engineering of the

University of Illinois at Urbana-Champaign

Apple v. Caltech

IPR2017-00211

Replacement - Apple 1102

HUGHES00883604

PROCEEDINGS

THIRTY-SEVENTH ANNUAL ALLERTON CONFERENCE ON COMMUNICATION, CONTROL, AND COMPUTING



Bruce Hajek R.S. Sreenivas Conference Co-Chairs



Conference held September 22, September 23, and September 24, 1999 Allerton House Monticello, Illinois

Sponsored by
The Coordinated Science Laboratory
and
The Department of Electrical and Computer Engineering
of the
UNIVERSITY OF ILLINOIS
at
Urbana-Champaign

TABLE OF CONTENTS

FORWORD		••••
-	WORKS I S.P. Meyn and R. Srikant (University of Illinois at Urbana-Champaign) S.P. Meyn (University of Illinois at Urbana-Champaign)	
REPRESENTATION AND F. Baccelli, S. Gau	EXPANSION OF (MAX, PLUS) LYAPUNOV EXPONENTSbert, and D. Hong	1
MARTINGALE PROBLEM T.G. Kurtz and R.F	IS AND LINEAR PROGRAMS FOR SINGULAR CONTROLI. Stockbridge	. 11
STATIONARY REFLECTE T. Konstantopoulos	ED LÉVY PROCESSES IN STOCHASTIC NETWORKSs and G. Last	21
ON THE IMPACT OF VAR Y. Joo, V. Ribeiro,	NABILITY ON THE BUFFER DYNAMICS IN IP NETWORKS	30
QUEUEING NETWORKS M. Armony and N.	WITH INTERACTING SERVICE RESOURCESBambos	42
Organizers:	: DECODING AND CHANNELS R. Koetter and R.E. Blahut (University of Illinois at Urbana-Champaign) A. Vardy (University of California, San Diego)	
A NEW UPPER BOUND O	N THE RELIABILITY FUNCTION OF THE GAUSSIAN CHANNEL	52
RECURSIVE DECODING (I. Dumer	OF REED-MULLER CODES	61
LOSSLESS COMPRESSION J.L. Fan, B. Marcus	N IN CONSTRAINED CODING	70
I-C: HYBRID/DISCRETE Chair:	-EVENT-DYNAMIC SYSTEMS R.S. Sreenivas (University of Illinois at Urbana-Champaign)	
MODELLING OF TIMED I R.S. Minhas and W	DISCRETE EVENT SYSTEMS	15
INTERACTING DISCRETE S. Abdelwahed and	E EVENT SYSTEMS	35
STABILITY ANALYSIS FO	OR INTERCONNECTED HYBRID SYSTEMS [. Ushio	13

DECENTRALIZED SUPERVISORY CONTROL OF CONCURRENT DISCRETE EVENT SYSTEMS WITH PARTIAL OBSERVATIONSS. Jiang and R. Kumar	. 103
A NEW PROBABILISTIC APPROACH TO CONGESTION CONTROL IN COMMUNICATION NETWORKS H. Mortazavian and J. Mirkovic	
A BIGRAPH MATCHING THEOREMS. Ayyorgun and R.L. Cruz	. 124
Organizer/Chair: Y. Shavitt (Bell Labs, Lucent Technologies)	
CHUNKS IN PLAN: LANGUAGE SUPPORT FOR PROGRAMS AS PACKETS	. 127
ON THE INTERFACE OF PROGRAMMABLE NETWORK ELEMENTSG. Hjálmtýsson	. 137
BOWMAN AND CANES: IMPLEMENTATION OF AN ACTIVE NETWORK	. 147
DESIGN OF A FLEXIBLE OPEN PLATFORM FOR HIGH PERFORMANCE ACTIVE NETWORKSS. Choi, D. Decasper, J. Dehart, R. Keller, J. Lockwood, J. Turner, and T. Wolf	157
IMPLEMENTING A CONCAST SERVICE	166
ACTIVE DISTRIBUTED MANAGEMENT FOR IP NETWORKS	176
I-F: SPACE-TIME METHODS FOR COMMUNICATION Chair: D. Sarwate	
(University of Illinois at Urbana-Champaign) NEW APPROACH FOR SPACE-TIME TRANSMITTER/RECEIVER DESIGN	186
INTERFERENCE SUPPRESSION FOR CDMA VIA A SPACE-TIME POWER MINIMIZATION BASED PREPROCESSOR WITH APPLICATIONS TO GPS W.L. Myrick, M.D. Zoltowski, and J.S. Goldstein	
SOFT-WEIGHTED TRANSMIT DIVERSITY FOR WCDMA	204
MULTIUSER DETECTION TECHNIQUES FOR COMBINED ARRAY PROCESSING AND SPACE-TIME BLOCK CODING B. Lu and X. Wang	214
A TRANSMIT ADAPTIVE ANTENNA SCHEME WITH FEEDBACK FOR WIRELESS COMMUNICATIONS	210

II-A: CODING THEORY II: ITERATIVE DECODING AND TURBO CODES Organizers: R. Koetter and R.E. Blahut (University of Illinois at Urbana-Champaign) Chair: R. Koetter
(University of Illinois at Urbana-Champaign)
CONCENTRATE221 T. Richardson and R. Urbanke
EFFICIENT ENCODING OF LOW-DENSITY PARITY-CHECK CODES
IRREGULAR TURBOCODES241 B.J. Frey and D.J.C. MacKay
ON QUASI-CYCLIC REPEAT-ACCUMULATE CODES
THE SERIAL CONCATENATION OF RATE-1 CODES THROUGH UNIFORM RANDOM INTERLEAVERS 260 H.D. Pfister and P.H. Siegel
II-B: STOCHASTIC NETWORKS II Organizers: S.P. Meyn and R. Srikant (University of Illinois at Urbana-Champaign) Chair: R. Srikant (University of Illinois at Urbana-Champaign)
QUEUE LENGTH ASYMPTOTICS FOR MARKOVIAN SERVICE NETWORKS
EXACT ASYMPTOTICS FOR 1-LIMITED EXPONENTIAL POLLING MODELS
INVARIANT RATE FUNCTIONS FOR DISCRETE TIME QUEUES
LARGE DEVIATIONS AND OPTIMALITY OF THE LARGEST WEIGHTED DELAY FIRST DISCIPLINE
ON ESTIMATING BUFFER OVERFLOW PROBABILITIES UNDER MARKOV-MODULATED INPUTS
INDUCED BURSTINESS IN GENERALIZED PROCESSOR SHARING QUEUES WITH LONG-TAILED TRAFFIC FLOWS
THE ASYMPTOTICS OF SELECTING THE SHORTEST OF TWO, IMPROVED

II-C: LEARNING ALGORITHMS IN SIGNAL PROCESSING Organizers: A. Singer and M. Feder (University of Illinois at Urbana-Champaign and Tel Aviv University) Chair: M. Feder (Tel Aviv University)	
UNIVERSAL FILTERING AND PREDICTION OF INDIVIDUAL SEQUENCES CORRUPTED BY NOISE A. Baruch and N. Merhav	
FAST RLS LAGUERRE ADAPTIVE FILTERING	
MACHINE LEARNING APPLICATIONS IN GRID COMPUTING	348
REDUNDANCY OF THE LEMPEL-ZIV CODES	358
THE INFORMATION BOTTLENECK METHOD	368
THEORY MEETS PRACTICE: UNIVERSAL SOURCE CODING WITH THE BURROWS WHEELER TRANSFORM	378
II-D: OPTICAL NETWORKS I Organizers: M. Médard and E. Modiano (University of Illinois at Urbana-Champaign and Massachusetts Institute of Technology) Chair: E. Modiano (Massachusetts Institute of Technology)	
OPTICAL SPACE COMMUNICATIONS AND NETWORKINGV.W.S. Chan	388
ON THE BENEFITS OF CONFIGURABILITY IN WDM NETWORKS E. Modiano and A. Narula-Tam	390
NONBLOCKING WDM NETWORKS WITH FIXED-TUNED TRANSMITTERS AND TUNABLE RECEIVERST. Lin and G. Sasaki	400
ON NEW ARCHITECTURES FOR WDM NETWORKS	402
ALL-OPTICAL LABEL SWAPPING WITH WAVELENGTH CONVERSION FOR WDM-IP NETWORKS WITH SUBCARRIER MULTIPLEXED ADDRESSING	414

II-E: COMMUNICATION SYSTEMS AND SERVICES

Chair: S. Lin

(University of Hawaii)

ON DISCRETE SUFFICIENT STATISTICS FOR ACQUISITION IN ASYNCHRONOUS
BAND-LIMITED CDMA SYSTEMS424 A. Mantravadi and V.V. Veeravalli
FREQUENCY SYNCHRONIZATION ALGORITHM FOR FREQUENCY HOPPING SYSTEM BASED ON SINGULAR VALUE DECOMPOSITION434
A. Pouttu
A SOFTWARE-ORIENTED STREAM CIPHER FOR CELLULAR AND PERSONAL COMMUNICATIONS SERVICES443
M. Zhang, A. Chan, and C. Carroll
BINARY RANK CRITERIA FOR PSK MODULATED SPACE-TIME CODES
TURBO CODES WITH ORTHOGONAL MODULATION IN DS-CDMA MOBILE RADIO SYSTEM WITH SHORT FRAME TRANSMISSION451
G. Li and Y.L. Guan
AN INTERACTIVE CONCATENATED TURBO CODING SYSTEM461 Y. Liu, H. Tang, S. Lin, and M.P.C. Fossorier
BI-DIRECTIONAL SOVA DECODING FOR TURBO-CODES
II-F: FADING CHANNELS AND POWER CONTROL Chair: D. Sarwate (University of Illinois at Urbana-Champaign)
ANALYSIS OF AN UP/DOWN POWER CONTROL ALGORITHM IN CDMA REVERSE LINK
UNDER FADING CONDITION L. Song, N. Mandayam, and Z. Gajic
A CLASS OF DISTRIBUTED ASYNCHRONOUS POWER CONTROL ALGORITHMS FOR CELLULAR
WIRELESS SYSTEMS
DISTRIBUTED CONNECTION ADMISSION CONTROL FOR POWER-CONTROLLED CELLULAR WIRELESS SYSTEMS
M. Xiao, N.B. Shroff, and E.K.P. Chong
INTERFERENCE AVOIDANCE AND DISPERSIVE CHANNELS: A NEW LOOK AT MULTICARRIER MODULATION
D.C. Popescu and C. Rose
PERFORMANCE OF OPTIMAL CODES ON GAUSSIAN AND RAYLEIGH FADING CHANNELS: A GEOMETRICAL APPROACH
S. Vialle and J. Boutros
IMPROVED MARKOV MODELS FOR FADING CHANNELS: ANALYSIS AND DESIGN

Organizers:	R. Koetter and R.E. Blahut (University of Illinois at Urbana-Champaign) N. Boston (University of Illinois at Urbana-Champaign)	
ON THE CLASSIFICATIO P. Gaborit, W.C. H	N OF EXTREMAL ADDITIVE CODES OVER GF(4)uffman, JL. Kim, and V. Pless	535
TWO FAST ALGORITHM GL. Feng	S IN THE SUDAN DECODING PROCEDURE	545
FROM WEIGHT ENUMER I. Duursma	RATORS TO ZETA FUNCTIONS	555
ALTERNATIVE APPROAGEODES	CHES TO THE COMPUTATION OF ERROR VALUES FOR HERMITIAN	557
_	FWORKS III S.P. Meyn and R. Srikant (University of Illinois at Urbana-Champaign) S.P. Meyn (University of Illinois at Urbana-Champaign)	
LONG-TAILED INPUTS	YMPTOTICS IN HOL SERVICE SYSTEMS WITH HETEROGENEOUS Likhanov and R.R. Mazumdar	567
SCHEDULING AND CON G. Weiss	TROL OF MANUFACTURING SYSTEMS — A FLUID APPROACH	577
MULTICLASS NETWORK POLICIES	S IN HEAVY TRAFFIC: ASYMPTOTIC OPTIMALITY OF TRACKING	587
SCHEDULING OPEN QUI S. Kumar	EUEING NETWORKS WITH SUFFICIENTLY FLEXIBLE RESOURCES	597
OPTIMALLY STABILIZIN P. Dupuis and R. A	NG CONTROLS FOR A DETERMINISTIC NETWORK MODELtar	607
III-C: ROBUST CONTRO Chair:	OL AND DECISION MAKING C. Beck (University of Illinois at Urbana-Champaign)	
PROBLEM IN THE PRESE	NT CONTROLLERS ON A BENCHMARK STRUCTURE PROTECTION ENCE OF CLASSIC EARTHQUAKES	617
RISK-SENSITIVE DECISI M.A. Shayman and	ON-THEORETIC TROUBLESHOOTINGI E. Fernández-Gaucherand	627

H ₂ CONTROL FOR MIXED DISTURBANCE REJECTION	
SOLVING POLYNOMIAL SYSTEMS IN ROBUST STABILITY ANALYSIS	
ROBUST, NEAR TIME-OPTIMAL CONTROL OF THIRD-ORDER UNCERTAIN SYSTEMSK.H. You and E.B. Lee	651
A NEW CONVEX RELAXATION FOR ROBUST H ₂ PERFORMANCE ANALYSIS OF UNCERTAIN LINEAR SYSTEMSL. El Ghaoui and E. Feron	655
A NEW RESULT ON THE BELLMAN EQUATION FOR EXIT TIME CONTROL PROBLEMS WITH CRITICAL GROWTH DYNAMICS M. Malisoff	657
III-D: OPTICAL NETWORKS II Organizers: M. Médard and E. Modiano (University of Illinois at Urbana-Champaign and Massachusetts Institute of Technology) Chair: M. Médard (University of Illinois at Urbana-Champaign)	
FAULT PROTECTION IN WDM MESH NETWORKS	
A COMPARISON OF ALLOCATION POLICIES IN WAVELENGTH ROUTING NETWORKSY. Zhu, G.N. Rouskas, and H.G. Perros	
OPTICAL BUFFERS FOR MULTI-TERABIT IP ROUTERS	
THE λ-SCHEDULER: A MULTIWAVELENGTH SCHEDULING SWITCH	
ON DIFFERENT ROUTING STRATEGIES IN TRANSPARENT ALL-OPTICAL NETWORKSO.K. Tonguz	699
III-E: COMMUNICATION NETWORKS Chair: R. Cruz (University of California, San Diego)	
TRANSMISSION POLICIES FOR TIME VARYING CHANNELS WITH AVERAGE DELAY CONSTRAINTS B.E. Collins and R.L. Cruz	709
FAIR ALLOCATION OF UTILITIES IN MULTIRATE MULTICAST NETWORKS	
ON THE USE OF MULTIPLE WORKING POINTS IN MULTICHANNEL ALOHA WITH DEADLINES D. Baron and Y. Birk	728

SPECIFICATION AND ANALYSIS OF A RELIABLE BROADCASTING PROTOCOL IN MAUDE	738
G. Denker, J.J. Garcia-Luna-Aceves, J. Meseguer, P.C. Ölveczky, J. Raju, B. Smith, and C.L. Talcott	730
MODELING AND ANALYSIS OF ACTIVE MESSAGES IN VOLATILE NETWORKS	748
IMPLEMENTATION OF AN ACTIVE CONGESTION CONTROL SCHEME IN NARROWBAND ATM NETWORKS	758
S. Sheth, J. Evans, A. Kulkarni, and G. Minden	
CPU SCHEDULING FOR ACTIVE PROCESSING USING FEEDBACK DEFICIT ROUND ROBINT. Wolf and D. Decasper	768
III-F: WIRELESS COMMUNICATION I: DETECTION AND ESTIMATION	
Organizers: V.V. Veeravalli and U. Madhow (Cornell University and University of Illinois at Urbana-Champaign)	
Chair: V.V. Veeravalli (Cornell University)	
PRECODING FOR SCATTERING FUNCTION ESTIMATION OF MOBILE CHANNELS USING OUTPUT CORRELATIONS ONLY	770
CORRELATIONS ONLY	.,,,
TWO-STAGE HYBRID ACQUISITION OF MULTICARRIER DIRECT-SEQUENCE	790
SPREAD-SPECTRUM SIGNALS	. 780
TRAINING SEQUENCE-BASED MULTIUSER CHANNEL ESTIMATION FOR	700
BLOCK-SYNCHRONOUS CDMA	. /90
SUPPRESSION OF HIGH-DENSITY, DYNAMIC NARROWBAND INTERFERENCE IN DS/CDMA	800
SPREAD-SPECTRUM SYSTEM	.000
LARGE SYSTEM PERFORMANCE OF REDUCED-RANK LINEAR FILTERS	.810
NONLINEAR MULTIUSER RECEIVERS WITH DISTRIBUTED POWER CONTROL IN CELLULAR RADIO NETWORKS	. 820
M K Varanasi	

Organizers:	IUNICATIONS II: SYSTEM CONSIDERATIONS IN PHYSICAL LAYER DESIGN.V. Veeravalli and U. Madhow (Cornell University and University of Illinois at Urbana-Champaign) U. Madhow (University of Illinois at Urbana-Champaign)	3N
THE CODING-SPREADING V.V. Veeravalli	G TRADEOFF IN CDMA SYSTEMS	831
SPECTRAL EFFICIENCY (B.M. Zaidel, S. Sha	DF RANDOMLY SPREAD DS-CDMA IN A MULTI-CELL MODEL mai, and S. Verdú	841
CDMA DESIGN THROUGH E. Biglieri, G. Caire	H ASYMPTOTIC ANALYSIS: FADING CHANNELS	851
PACKING SPHERES IN THE THE NON-COHERENT MUL. Zheng and D.N.C	IE GRASSMANN MANIFOLD: A GEOMETRIC APPROACH TO JLTI-ANTENNA CHANNEL	861
BLIND ADAPTIVE MULTI GENERALIZED RANDOM J.H. Cho and J.S. Le	USER DETECTION FOR DS/SSMA COMMUNICATIONS WITH SPREADING IN A FREQUENCY-SELECTIVE FADING CHANNEL	371
TRAFFIC AIDED MULTIUS NETWORKSB. Chen and L. Tong	SER DETECTION FOR PACKET SWITCHING RANDOM ACCESS/CDMA 8	381
	IV R.E. Blahut (University of Illinois at Urbana-Champaign)	
OPTICAL CHANNELS WIT A. Khandekar and R	H DEAD TIME	91
AWGN CODING THEOREM H. Jin and R.J. McEl	4S FOR SERIAL TURBO CODES8	93
LINEAR CODES OVER Z/(2 J.A. Wood	OF CONSTANT EUCLIDEAN WEIGHT8	95
SPACE-TIME TURBO COD Y. Liu and M.P. Fitz	ES8	97
K.R. Narayanan	NCATENATED SPACE-TIME CODES89	
SOFT OUTPUT AND ITERA R. Sivasankaran and	TIVE STACK DECODING90 S.W. McLaughlin)1
ALGEBRAIC GEOMETRIC GILBERT-VARSHAMOV B H. Maharaj	CODES AND AN IMPROVEMENTS ON THE OUND90)3

IV-C: STOCHASTIC SYSTEMS AND CONTROL Chair: G. Dullerud (University of Illinois at Urbana-Champaign)
SINGLE-SAMPLE-PATH-BASED OPTIMIZATION OF MARKOV DECISION PROCESS
TRACKING CAPABILITY ANALYSIS OF THE LMS ALGORITHM FOR FIR SYSTEMS WITH AR COEFFICIENTS
ADAPTIVE OPTIMAL PREDICTION FOR MIMO STOCHASTIC SYSTEMS USING CANONICAL FORMS 925 B. Shahrrava and J.D. Aplevich
SDRE ESTIMATOR PERFORMANCE IN A HARMONIC DISTORTION PROBLEM
STOCHASTIC DIFFERENCE EQUATIONS WITH TIME DELAYS
CONTROL UNDER COMMUNICATION CONSTRAINTS
IV-D: SPACE-TIME PROCESSING I Organizer/Chair: A. Nehorai (University of Illinois at Chicago)
SPACE-TIME FADING CHANNEL ESTIMATION IN UNKNOWN SPATIALLY CORRELATED NOISE948 A. Dogandžić and A. Nehorai
BLAST TRAINING: ESTIMATING CHANNEL CHARACTERISTICS FOR HIGH CAPACITY SPACE-TIME WIRELESS
ANALYSIS OF THE PARALLEL INTERFERENCE CANCELER FOR DS/CDMA SIGNALS967 R. Chandrasekaran and J.J. Shynk
BEARING ESTIMATION IN A RICEAN CHANNEL
OPTIMAL DOWNLINK BEAMFORMING USING SEMIDEFINITE OPTIMIZATION
AN OVERVIEW OF A SIMULATION ENVIRONMENT TO STUDY THE IMPACT OF NON-IDEAL HARDWARE ON ARRAY PROCESSING
SIGNAL REPRESENTATIONS FOR TRANSMIT-RECEIVE ANTENNA ARRAYS
BLOCK SPACE-TIME ANTENNA PRECODING/DECODING FOR GENERALIZED MULTICARRIER COMMUNICATIONS IN UNKNOWN MULTIPATH

LEAST-SQUARES MUL	TI-USER FREQUENCY-DOMAIN CHANNEL ESTIMATION ELESS COMMUNICATION SYSTEMS	1026
FOR BROADBAND WIR T.A. Thomas, F.V	W. Vook, and K.L. Baum	
IV-E: MANUFACTURIN Organizer/Chair	NG SYSTEMS : S. Reveliotis	
Organizen Chan	(Georgia Institute of Technology)	
SYSTEMS	PROCESS MODELING FOR CONTROL SWITCHING OF DISCRETE EVENT	036
H. Darabi and M.	A. Jafari	
SUPERVISORY CONTRO X. Guan and L.E.	OL OF CONTRADICTIONS IN HIERARCHICAL TASK CONTROLLERS I Holloway	042
IV-F: STOCHASTIC NE	TWORKS IV S.P. Meyn and R. Srikant	
_	(University of Illinois at Urbana-Champaign)	
Chair:	R. Srikant (University of Illinois at Urbana-Champaign)	
MULTIUSER RECEIVERS D.N.C. Tse	S, RANDOM MATRICES AND FREE PROBABILITY10)55
STABILITY PROPERTIES R. Vijayakumar an	S OF INCREMENTAL REDUNDANCY IN CDMA PACKET DATA NETWORKS 10 Id K.M. Wasserman)65
P. Marbach	SSES IN A DIFFERENTIATED SERVICES NETWORK10	
COMPARING TANDEM C E. Altman, G. Koo	QUEUEING SYSTEMS AND THEIR FLUID LIMITS10 le, and T. Jiménez	85
WAITING TIME ASYMPT	OTICS FOR TIME VARYING MULTISERVER QUEUES	95
WITH ABANDONMENT A. Mandelbaum, W.	ND RETRIALS	,,
	R SMALL BUFFERS: AN INSENSITIVITY RESULT110	05
OVERFLOW AND LOSSE: B. Tsybakov and N	S IN A NETWORK QUEUE WITH SELF-SIMILAR INPUT11 D. Georganas	13
V-A: STOCHASTIC NET	WORKS V	
	S.P. Meyn and R. Srikant (University of Illinois at Urbana-Champaign)	
Chair:	S.P. Meyn (University of Illinois at Urbana-Champaign)	
CHOKE A STATELESS MI	ECHANISM FOR PROVIDING QUALITY OF SERVICE	
IN THE INTERNET		:2
N. Fall alid D. Fladi	****	

OPTIMAL ROUTING TO A E. Altman, S. Bhula	M PARALLEL QUEUES WITH NO BUFFERSai, B. Gaujal, and A. Hordijk	1132
THE EFFECT OF SCALE O M. Siler and J. Wal	ON INTERNET QUALITYrand	1142
	NE INPUT-BUFFERED SWITCHES WITH MARKOV-MODULATED J. Walrand	1152
Organizers:	UNICATIONS III: NETWORKING ISSUES V.V. Veeravalli and U. Madhow (Cornell University and University of Illinois at Urbana-Champaign) V.V. Veeravalli (University of Illinois at Urbana-Champaign)	
	ON FOR SPREAD-SPECTRUM COMMUNICATIONS OVER MULTIPATH	1162
ADAPTIVE REDUNDANC T. Ji and W.E. Star	CY RETRANSMISSION PROTOCOLS FOR WIRELESS NETWORKSk	1171
RESOURCE POOLING AN S.V. Hanly and D.N	ND EFFECTIVE BANDWIDTHS FOR CDMA ANTENNA ARRAYS	1181
ROUTING FOR MAXIMU JH. Chang and L.	M SYSTEM LIFETIME IN WIRELESS AD-HOC NETWORKSTassiulas	1191
	RELESS SENSOR NETWORK	1201
V-C: MULTIUSER DETE Chair:	CCTION M.K. Varanasi (University of Colorado at Boulder)	
BLIND ADAPTIVE NONC D. Das and M.K. V	COHERENT MULTIUSER DETECTION FOR NONLINEAR MODULATION	1211
FOR OVERSATURATED	I-COHERENT NEAR-OPTIMAL MULTIUSER DETECTION MA COMMUNICATION	1221
CELLULAR DS-CDMA	DECISION FEEDBACK FOR ASYNCHRONOUS odward, and M.L. Honig	1236
MULTIUSER EQUALIZA' WITH AND WITHOUT DI R.R. Müller	TION FOR RANDOM SPREADING: LIMITS OF DECORRELATION ECISION-FEEDBACK	1246

V-D: SPACE-TIME PROCESSING II Organizer/Chair: A. Nehorai (University of Illinois at Chicago)
SCHEDULING OF SWITCHED MULTIBEAM ANTENNAS IN A MULTIPLE ACCESS ENVIRONMENT 1256 A. Logothetis and H.V. Poor
MULTIPLE ANTENNA DIFFERENTIAL MODULATION
SPACE-TIME ZERO FORCING EQUALIZATION FOR 3G CDMA FORWARD LINK TO RESTORE ORTHOGONALITY OF CHANNEL CODES1274 M.D. Zoltowski and T.P. Krauss
A NOVEL SPACE-TIME SPREADING SCHEME FOR WIRELESS CDMA SYSTEMS
V-E: CODING FOR MAGNETIC CHANNELS Organizer/Chairs: E. Kurtas (Quantrum Corporation)
2D EQUALIZATION FOR PAGE-ORIENTED DATA STORAGE SYSTEMS
TOWARDS SOFT OUTPUT APP DECODING FOR NONSYSTEMATIC NONLINEAR BLOCK CODES 1304 K.D. Anim-Appiah and S.W. McLaughlin
LOW DENSITY PARITY CHECK CODES FOR MAGNETIC RECORDING
DESIGN CONSIDERATIONS FOR CONCATENATING CONVOLUTIONAL CODES WITH PARTIAL RESPONSE CHANNELS
TURBO CODES FOR TWO-TRACK MAGNETIC RECORDING SYSTEMS
LIST OF AUTHORS

Irregular Turbocodes

Brendan J. Frey

Computer Science, University of Waterloo Electrical and Computer Engineering, University of Illinois at Urbana http://www.cs.uwaterloo.ca/~frey

David J. C. MacKay

Department of Physics, Cavendish Laboratories Cambridge University http://wol.ra.phy.cam.ac.uk/mackay

Abstract

Recently, several groups have increased the coding gain of iteratively decoded Gallager codes (low density parity check codes) by varying the number of parity check equations in which each codeword bit participates. In regular turbocodes, each "systematic bit" participates in exactly 2 trellis sections. We construct irregular turbocodes with systematic bits that participate in varying numbers of trellis sections. These codes can be decoded by the iterative application of the sum-product algorithm (a low-complexity, more general form of the turbodecoding algorithm). By making the original rate 1/2 turbocode of Berrou et al. slightly irregular, we obtain a coding gain of 0.15 dB at a block length of N=131,072, bringing the irregular turbocode within 0.3 dB of capacity. Just like regular turbocodes, irregular turbocodes are linear-time encodable.

1 Introduction

Recent work on irregular Gallager codes (low density parity check codes) has shown that by making the codeword bits participate in varying numbers of parity check equations, significant coding gains can be achieved [1–3]. Although Gallager codes have been shown to perform better than turbocodes at BERs below 10^{-5} [4]¹, until recently Gallager codes performed over 0.5 dB worse than turbocodes for BERs greater than 10^{-5} . However, in [3], Richardson *et al.* found irregular Gallager codes that perform 0.16 dB *better* than the original turbocode at BERs greater than 10^{-5} [5] for a block length of $N \approx 131,072$.

¹Gallager codes to not exhibit decoding errors, only decoding failures, at long block lengths with N > 5,000.

In this paper, we show that by tweaking a turbocode so that it is irregular, we obtain a coding gain of 0.15 dB for a block length of N=131,072. For example, an N=131,072 irregular turbocode achieves $E_b/N_0=0.48$ dB at BER = 10^{-4} , a performance similar to the best irregular Gallager code published to date [3]. By further optimizing the degree profile, the permuter and the trellis polynomials, we expect to find even better irregular turbocodes. Like their regular cousins, irregular turbocodes exhibit a BER flattening due to low-weight codewords.

2 Irregular turbocodes

In Fig. 1, we show how to view a turbocode so that it can be made irregular. The first picture shows the set of systematic bits (middle row of discs) being fed directly into one convolutional code (the chain at the top) and being permuted before being fed into another convolutional code (the chain at the bottom). For a rate 1/2 turbocode, each constituent convolutional code should be rate 2/3 (which may, for example, be obtained by puncturing a lower-rate convolutional code).

Since the order of the systematic bits is irrelevant, we may also introduce a permuter before the upper convolutional code, as shown in the second picture. In the third picture, we have simply drawn the two permuters and convolutional codes side by side.

For long turbocodes, the values of the initial state and the final state of the convolutional chains do not significantly influence performance (e.g., see [6]). So, as shown in the fourth picture, we can view a turbocode as a code that copies the systematic bits, permutes both sets of these bits, and then feeds them into a convolutional code. We refer to this turbocode as "regular", since each systematic bit is copied exactly once.

The final picture illustrates one way the above turbocode can be made irregular. Some of the systematic bits are "tied" together, in effect causing some systematic bits to be replicated more than once. Notice that too keep the rate of the overall code fixed at 1/2, some extra parity bits must be punctured.

More generally, an irregular turbocode has the form shown in Fig. 2, which is a type of "trellis-constrained code" as described in [7]. We specify a degree profile, $f_d \in [0, 1], d \in \{1, 2, \ldots, D\}$. f_d is the fraction of codeword bits that have degree d and D is the maximum degree. Each codeword bit with degree d is repeated d times before being fed into the permuter. Several classes of permuter lead to linear-time encodable codes. In particular, if the bits in the convolutional code are partitioned into "systematic bits" and "parity bits", then by connecting each parity bit to a degree 1 codeword bit, we can encode in linear time.

The average codeword bit degree is

$$\bar{d} = \sum_{d=1}^{D} d \cdot f_d \tag{1}$$

The overall rate R of an irregular turbocode is related to the rate R' of the convolutional code and the average degree \bar{d} by

$$\bar{d}(1 - R') = 1 - R. \tag{2}$$

So, if the average degree is increased, the rate of the convolutional code must also be increased to keep the overall rate constant. This can be done by puncturing the convolutional code or by designing a new, higher rate convolutional code.

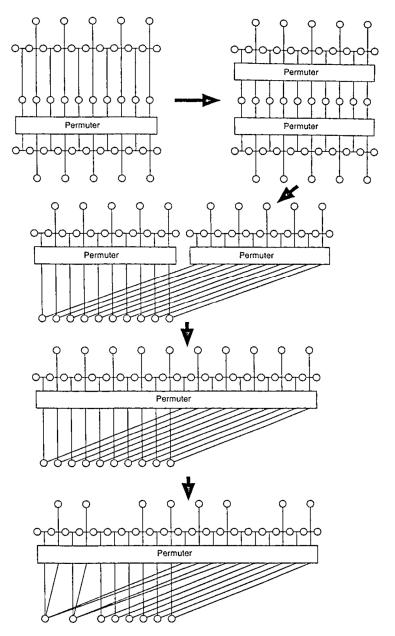


Figure 1: The first 4 pictures show that a turbocode can be viewed as a code that copies the systematic bits, permutes both sets of these bits and then feeds them into a convolutional code. The 5th picture shows how a turbocode can be made irregular by "tying" some of the systematic bits together, *i.e.*, by having some systematic bits replicated more than once. Too keep the rate fixed, some extra parity bits must be punctured. Too keep the block length fixed, we must start with a longer turbocode.

3 Decoding irregular turbocodes

Fig. 2 can be interpreted as the graphical model [6, 8–10] for the irregular turbocode. Decoding consists of the iterative application of the sum-product algorithm (a low-complexity, more general form of turbodecoding) in this graph.

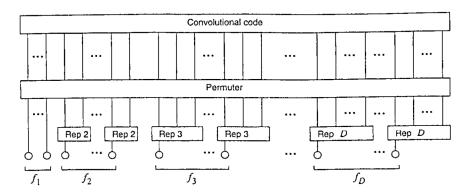


Figure 2: A general *irregular turbocode*. For d = 1, ..., D, fraction f_d of the codeword bits are repeated d times, permuted and connected to a convolutional code.

After receiving the channel output, the decoder computes the channel output \log -likelihood ratios for the N codeword bits,

$$L_1^0, L_2^0, \dots, L_N^0,$$
 (3)

and then repeats each log-likelihood ratio appropriately. If codeword bit i has degree d, we set

$$L_{i,1} \leftarrow L_i^0, \ L_{i,2} \leftarrow L_i^0, \dots, \ L_{i,d} \leftarrow L_i^0.$$
 (4)

Next, the log-likelihood ratios are permuted and fed into the BCJR algorithm [11] for the convolutional code. The BCJR algorithm assumes the inputs are a priori log-probability ratios and uses the forward-backward algorithm [12] to compute a set of a posteriori log-probability ratios. If codeword bit i has degree d, the algorithm produces d a posteriori log-probability ratios,

$$L'_{i,1}, L'_{i,2}, \dots, L'_{i,d}.$$
 (5)

For a regular turbocode, there are just two *a posteriori* log-probability ratios, $L'_{i,1}$ and $L'_{i,2}$, for each degree 2 bit and they correspond to the "extrinsic information" produced by each constituent convolutional code.

The current estimate of the log-probability ratio for bit i given the channel output is

$$\hat{L}_i \leftarrow L_i^0 + \sum_{k=1}^d (L'_{i,k} - L_{i,k}).$$
 (6)

To compute the inputs to the BCJR algorithm needed for the next iteration, we subtract off the corresponding outputs from the BCJR algorithm produced by the previous iteration:

$$L_{i,k} \leftarrow \hat{L}_i - L'_{i,k}. \tag{7}$$

So, each iteration consists of computing the inputs to the BCJR algorithm, permuting the inputs, applying the BCJR algorithm, permuting the outputs of the BCJR algorithm, and taking the repetitions into account to combine the outputs to form estimates of the log-probability ratios of the codeword bits given the channel output.

In our simulations, after each iteration, we check to see if the current decision gives a codeword. If it does, the iterations terminate and otherwise, the decoder iterates further until some maximum number of iterations is reached and a decoding failure is declared.

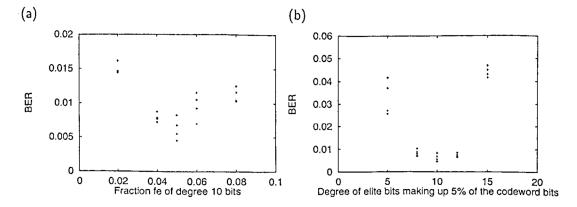


Figure 3: (a) shows the effect of changing the fraction of elite bits on the BER, while keeping the degree of the elite bits fixed at 10. (b) shows the effect of changing the degree of the elite bits on the BER, while keeping the fraction of elite bits fixed at 0.05. For each fraction and degree, the performance of 4 randomly drawn permuters is shown.

4 Selecting the profile

Finding a good profile is not trivial, since the best profile will depend on the parameters of the convolutional code, the permuter and the distortion measure (bit error rate, block error rate, decoding failure rate, high-weight decoding failure rate, etc.)

The results we report in this paper were obtained by making small changes to a block length N=10,000 version of the original rate R=1/2 turbocode proposed by Berrou et al.. In this turbocode, $f_1=f_2=1/2$ (see Fig. 2) and the convolutional code polynomials are 37 and 21 (octal). The taps associated with polynomial 37 are connected to the degree 2 codeword bits, 1/2 of the taps associated with polynomial 21 are connected to the degree 1 bits, and the remaining 1/2 of the taps associated with polynomial 21 are punctured, giving the required convolutional code rate of R'=2/3.

To simplify our search, we considered profiles where besides degrees 1 and 2, only one other degree, e for "elite", had a nonzero fraction. So, for a code with overall rate R and fraction f_e of degree e elite bits, we have

$$f_1 = 1 - R = 1/2,$$

$$f_2 = 1 - f_1 - f_e = 1/2 - f_e.$$
 (8)

In this restricted class of codes, irregularity is governed by two parameters, e and f_e .

From (2) it is clear that when the average degree is increased, the rate of the convolutional code must also be increased to keep the overall rate at 1/2. We increased the rate of the punctured convolutional code by further puncturing the taps associated with polynomial 21 to obtain a convolutional code with rate

$$R' = 1 - \frac{1 - R}{\bar{d}} = 1 - \frac{1/2}{1/2 + 2(1/2 - f_e) + ef_e}.$$
 (9)

So, in the codes we explored, the level of puncturing was quite high and some extra low-weight codewords were introduced.

To begin with, we made an irregular turbocode with e = 10 (chosen using intuition) and varied f_e from 0.02 to 0.08 while measuring the BER at $E_b/N_0 = 0.6$ dB. In each

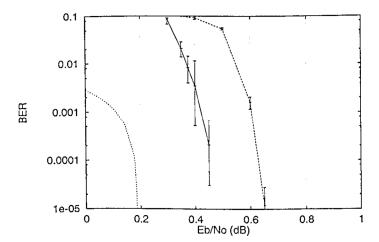


Figure 4: Performance of the original block length N=131,072 turbocode (dashed line) and one of its irregular cousins (solid line). These results are for irregular turbocodes obtained by tweaking the original turbocode – we are currently searching for optimal degree profiles, permuters and trellis polynomials.

experiment, we simulated enough blocks to obtain a relatively small confidence interval. The results are shown in Fig. 3a, which indicates that for degree 10 elite bits, the best fraction is roughly 0.05. Next, we kept the fraction of elite bits fixed at $f_e = 0.05$ and we varied the degree of the elite bits. The results are shown in Fig. 3b, which indicates that for a fraction of 0.05, the best degree is roughly 10.

These results show that for e = 10, $f_e = 0.05$ is a good fraction and that for $f_e = 0.05$, e = 10 is a good degree. However, values of e and f_e that give good profiles are probably correlated, so we are currently extending our search.

5 Results

Fig. 4 shows the simulated BER- E_b/N_0 curves for the original block length N=131,072 regular turbocode (dashed line) and its irregular cousin (solid line), using profile e=10, $f_e=0.05$.

The irregular turbocode clearly performs better than the regular turbocode for BER $> 10^{-4}$. At BER $= 10^{-4}$, the N = 131,072 irregular turbocode is 0.3 dB from capacity, a 0.15 dB improvement over the regular turbocode.

For high E_b/N_0 , most of the errors for the irregular turbocode were due to low-weight codewords. According to preliminary results, the distribution of error weights appears to indicate that the flattening effect for the particular N=131,072 irregular turbocode we constructed occurs at a higher BER than it does for the regular turbocode. However, the flattening effect is highly sensitive to the technique used to construct the permuter (we drew it at random) and the design of the convolutional code (we just further punctured the convolutional code used in the original turbocode). We are currently experimenting with techniques for lowering the level of the flattening effect (e.g., see [13]).

6 Conclusions

We have shown that by making the original, regular turbocode irregular, a coding gain of 0.15 dB is obtained, bringing the irregular turbocode within 0.3 dB of capacity at a BER of 10⁻⁴. This irregular turbocode performs in the same regime as the best known irregular Gallager code.

We emphasize that we obtained these results by tweaking the regular turbocode originally introduced by Berrou et al.. We believe we will be able to improve these performance curves significantly, both by exploring the polynomials used in the convolutional code and by adjusting the degree profile and the permuter structure. (One way to speed up the search is to extend the method of "density evolution" [3] to models with state.) In particular, we are investigating ways to select the permuter and the polynomials to eliminate low-weight codewords, thus reducing the flattening effect [13–15]. We are also studying ways of constraining the degree 1 "parity" bits (i.e., increasing their degree) to eliminate low-weight codewords.

References

- [1] M. G. Luby, M. Mitzenmacher, M. A. Shokrollahi, and D. A. Spielman, "Improved low-density parity-check codes using irregular graphs and belief propagation," in *Proceedings of IEEE International Symposium on Information Theory*, 1998.
- [2] D. J. C. MacKay, S. T. Wilson, and M. C. Davey, "Comparison of constructions of irregular Gallager codes," *IEEE Transactions on Communications*, vol. 47, October 1999.
- [3] T. Richardson, A. Shokrollahi, and R. Urbanke, "Design of provably good low density parity check codes." Submitted to *IEEE Transactions on Information Theory*, July 1999.
- [4] D. J. C. MacKay, "Good error-correcting codes based on very sparse matrices," *IEEE Transactions on Information Theory*, vol. 45, pp. 399-431, March 1999.
- [5] C. Berrou and A. Glavieux, "Near optimum error correcting coding and decoding: Turbo-codes," *IEEE Transactions on Communications*, vol. 44, pp. 1261–1271, October 1996.
- [6] B. J. Frey, Graphical Models for Machine Learning and Digital Communication. Cambridge MA.: MIT Press, 1998. See http://www.cs.utoronto.ca/~frey.
- [7] B. J. Frey and D. J. C. MacKay, "Trellis-constrained codes," in Proceedings of the 35th Allerton Conference on Communication, Control and Computing 1997, 1998. Available at http://www.cs.utoronto.ca/~frey.
- [8] N. Wiberg, H.-A. Loeliger, and R. Kötter, "Codes and iterative decoding on general graphs," European Transactions on Telecommunications, vol. 6, pp. 513-525, September/October 1995.
- [9] F. R. Kschischang and B. J. Frey, "Iterative decoding of compound codes by probability propagation in graphical models," *IEEE Journal on Selected Areas in Communications*, vol. 16, pp. 219–230, February 1998.

- [10] R. J. McEliece, D. J. C. MacKay, and J. F. Cheng, "Turbo-decoding as an instance of Pearl's 'belief propagation' algorithm," *IEEE Journal on Selected Areas in Com*munications, vol. 16, February 1998.
- [11] L. R. Bahl, J. Cocke, F. Jelinek, and J. Raviv, "Optimal decoding of linear codes for minimizing symbol error rate," *IEEE Transactions on Information Theory*, vol. 20, pp. 284-287, March 1974.
- [12] L. E. Baum and T. Petrie, "Statistical inference for probabilistic functions of finite state markov chains," *Annals of Mathematical Statistics*, vol. 37, pp. 1559-1563, 1966.
- [13] M. Öberg and P. H. Siegel, "Application of distance spectrum analysis to turbocode performance improvement," in Proceedings of the 35th Allerton Conference on Communication, Control and Computing 1997, 1998.
- [14] D. Divsalar and F. Pollara, "Turbo-codes for PCS applications," in *Proceedings of the International Conference on Communications*, pp. 54-59, 1995.
- [15] D. Divsalar and R. J. McEliece, "Effective free distance of turbocodes," *Electronics Letters*, vol. 32, pp. 445–446, February 1996.

